

Defⁿ (bounded): Let $A \subseteq \mathbb{R}$ and let $f: A \rightarrow \mathbb{R}$ and let c be a cluster point of A . We say f is bounded on a nbd of c if there exists a δ -nbd of c , $N_\delta(c)$ and a constant M s.t. $|f(x)| \leq M \quad \forall x \in A \cap N_\delta(c)$.

Theorem: Let $A \subseteq \mathbb{R}$ and $f: A \rightarrow \mathbb{R}$ has a limit at $c \in \mathbb{R}$ then f is bounded on some nbd of c .

Proof: If $l = \lim_{x \rightarrow c} f(x)$.

For given $\epsilon (=1)$, $\exists \delta > 0$ s.t.

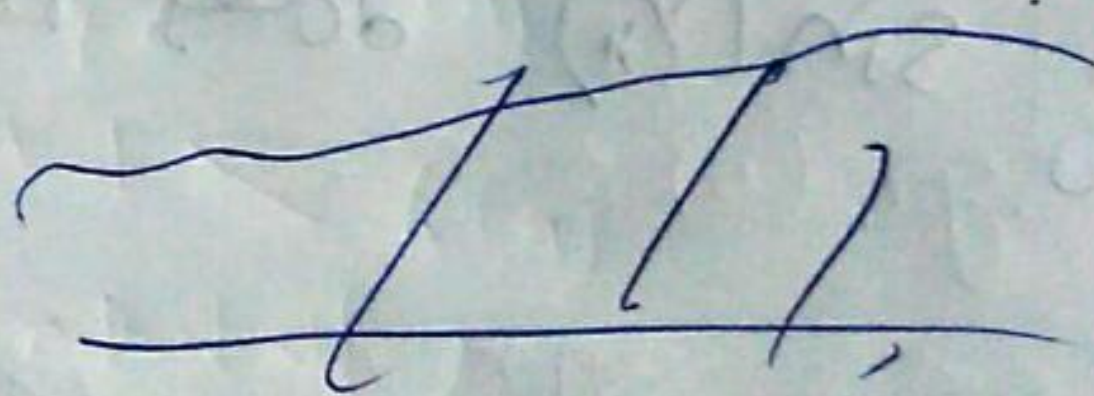
$$\text{If } 0 < |x - c| < \delta, \text{ then } |f(x) - l| < 1 \\ \Rightarrow |f(x)| < |l| + 1$$

$$\text{Here } M = |l| + 1$$

In case if $c \in A$, we can take

$$M = \sup \{ |f(c)|, |l| + 1 \}$$

then $|f(x)| \leq M$ for $x \in N_\delta(c) \cap A$.



Ex
Converse of this theorem may not be true i.e., f is a bdd funn on some nbd of c then $\lim_{x \rightarrow c} f(x)$ may or may not exist.

e.g., $f(x) = \text{sgn}(x) = \begin{cases} 1, & x > 0 \\ 0, & x = 0 \\ -1, & x < 0 \end{cases}$

Range of this funⁿ is ~~abdt~~ bounded set.
So, in the nbd of 0, it is basically a bounded set.

i.e., $|f(x)| \leq 1 \quad \forall x \in N_\delta(0)$

But $\lim_{x \rightarrow 0} \text{sgn}(x)$ does not exist.

Note:- If $f(x)$ is not bounded in some nbd of c , then $\lim_{x \rightarrow c} f(x)$ will not exist.

e.g., $\lim_{x \rightarrow 0} \frac{1}{x}$ does not exist.

$x_n = \frac{1}{n} \in N_{\frac{1}{n}}(0)$ but $f(x_n) = \frac{1}{x_n} = n$

$\downarrow$
 0 as $n \rightarrow \infty$ $\therefore \{n\}$ is a divergent sequence.

Since $\frac{1}{x}$ is unbounded in the nbd of 0, so $\lim_{x \rightarrow 0} \frac{1}{x}$ does not exist.

Defⁿ: Let $A \subseteq \mathbb{R}$ and let f and g be funⁿs defined on A to $\mathbb{R}$. Then we define

$$(f \pm g)(x) = f(x) \pm g(x) \quad \forall x \in A$$

$$(fg)(x) = f(x)g(x) \quad "$$

$$(bf)(x) = b f(x) \quad \text{where } b \in \mathbb{R}.$$

If $h(x) \neq 0$ for $x \in A$ then

$$\left(\frac{f}{h}\right)(x) = \frac{f(x)}{h(x)} \quad \forall x \in A.$$

Theorem: Let $A \subseteq \mathbb{R}$, let f and g are funⁿs from A to $\mathbb{R}$, and let $c \in \mathbb{R}$ be a cluster pt. of A . Further let $b \in \mathbb{R}$, then

(a) If $\lim_{n \rightarrow c} f = l$ i.e., $\lim_{n \rightarrow c} f(n) = l$ and $\lim_{n \rightarrow c} g = m$

then $\lim_{n \rightarrow c} (f \pm g) = l \pm m$

$\lim_{n \rightarrow c} (fg) = lm$ & $\lim_{n \rightarrow c} (bf) = bl$

(b) If $h: A \rightarrow \mathbb{R}$, if $h(n) \neq 0 \forall n \in A$, and if $\lim_{n \rightarrow c} h = H \neq 0$ then $\lim_{n \rightarrow c} \frac{f}{h} = \frac{L}{H}$.

Note:- Let $A \subseteq \mathbb{R}$. Let $f_1, \dots, f_m$ be funⁿs on A to $\mathbb{R}$ and having limiting values, say $l_1, l_2, \dots, l_m$ respectively, at c , cluster pt. of A . Then $\lim_{n \rightarrow c} (f_1 + f_2 + \dots + f_m) = l_1 + l_2 + \dots + l_m$.
 $\lim_{n \rightarrow c} (f_1 f_2 \dots f_m) = l_1 l_2 \dots l_m$

Ex: 1. Find $\lim_{n \rightarrow 2} \frac{n^3 - 4}{n^2 + 1} \leftarrow f(n)$

$h(n) \neq 0$ in the nbd of 2 $\leftarrow h(n)$
 $\lim_{n \rightarrow 2} \frac{n^3 - 4}{n^2 + 1} = \frac{2^3 - 4}{2^2 + 1} = \frac{4}{5}$

2. $\lim_{n \rightarrow 2} \frac{n^2 - 4}{3n - 6}$ ($\frac{0}{0}$ form)

$= \lim_{n \rightarrow 2} \frac{(n-2)(n+2)}{3(n-2)} = \lim_{n \rightarrow 2} \frac{n+2}{3} = \frac{4}{3}$

3. If p is a polynomial of degree n
 $p(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_0, x \in \mathbb{R}$
 $a_0, a_1, \dots, a_n \in \mathbb{R}$.

$$\text{Then } \lim_{x \rightarrow c} p(x) = a_n c^n + a_{n-1} c^{n-1} + \dots + a_0 \\ = p(c)$$

4. If $p(x)$ and $q(x)$ are polynomial functions on $\mathbb{R}$ and if $q(c) \neq 0$, then

$$\lim_{x \rightarrow c} \frac{p(x)}{q(x)} = \frac{p(c)}{q(c)}.$$