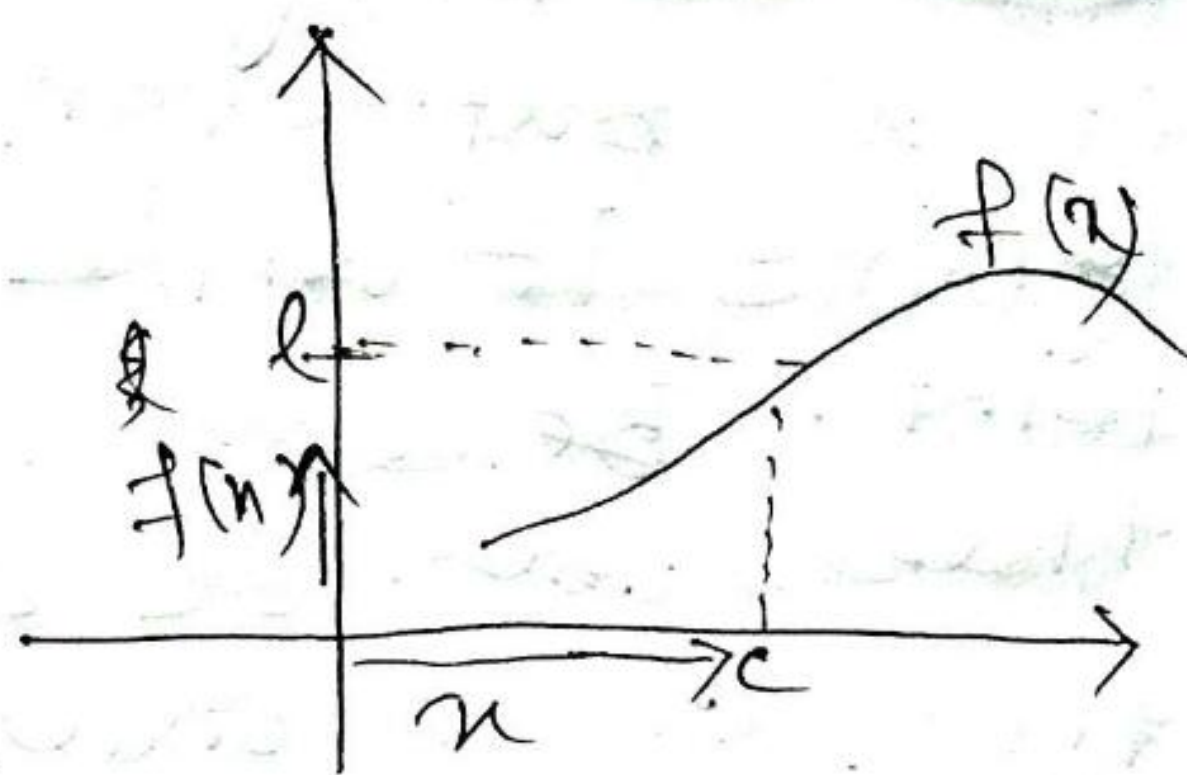


Limit of Functions

Let f is a real valued function. Let $c \in \mathbb{R}$.

$\lim_{x \rightarrow c} f(x) = l$ means $f(x)$ is close to l

when x is sufficiently close to c .



The important part or necessary part, in order to get the limit of a the funⁿ $f(x)$ at a point c —

(i) The funⁿ must be define at all points which is very very close to c . At c , $f(x)$ may or may not be defined.

(ii) ~~At~~ The point c may or may not be available in the set A (the domain of the funⁿ) still we can find the limiting value of $f(x)$ when x approaches to c .

Note:- The point c may or may not be in the set, but still $\lim_{x \rightarrow c} f(x)$ can be defined, provided $x \in N_\delta(c) \cap A$.

Defⁿ (cluster point of a set A):- Let $A \subseteq \mathbb{R}$. ($\neq \emptyset$)

A point $c \in \mathbb{R}$ is said to be a cluster point of A if for every $\delta > 0$ there exists at least one point $x \in A$, $x \neq c$ s.t. $|x - c| < \delta$.

OR.

A point c is a cluster point (limit point) of the set A if every δ -neighbourhood (nbd)

$$\leftarrow N_\delta(c) = (c - \delta, c + \delta)$$

$N_\delta(c)$, $(c-\delta, c+\delta)$ of c contains at least one point of A distinct from c .

Note:- c may or may not be a member of A , but even if it is in A , still we ignore it while deciding c as a cluster point. Because we need to show that there ~~exists~~ be a point in $N_\delta(c) \cap A$ distinct from c in order for c to be a cluster point of A .

e.g., $A = \{1, 2\}$

Obviously 1 is not the cluster point of A .

Because $N_{1/2}(1) \cap A = \{1\}$

So there does not exist ~~any~~ $x \in A$

any x in $N_{1/2}(1) \cap A$.

Hence 1 can't be cluster pt.

Similarly 2 can't be.

Theorem:- A number $c \in \mathbb{R}$ is a cluster point (limit point) of a subset A of $\mathbb{R}$ if and only if (iff) there exists a sequence $\{a_n\}$ in A such that $\lim_{n \rightarrow \infty} a_n = c$ and $a_n \neq c \forall n \in \mathbb{N}$.

Proof:- If c is a cluster point of A , then for any $n \in \mathbb{N}$, $N_{1/n}(c)$ contains at least one point a_n in A distinct from c .

Then $a_n \in A$, $a_n \neq c$ and $|a_n - c| < \frac{1}{n}$.

$$\Rightarrow \lim_{n \rightarrow \infty} a_n = c.$$

Conversely, given a sequence $\{a_n\}$ in A s.t. $\lim_{n \rightarrow \infty} a_n = c$.

For any $\varepsilon > 0$ $\exists k$ s.t. if $n \geq k$ then $a_n \in N_\varepsilon(c)$.

$\therefore N_\varepsilon(c)$ contains a_n , for $n \geq k$, which belongs to A and distinct from c .

$\Rightarrow c$ is the cluster point of A .

Ex:-

1. $A = (0, 1) \subseteq \mathbb{R}$.

Every point including 0 & 1 are cluster pt.

2. $A = \text{finite set of real numbers}$
 $= \{a_1, a_2, \dots, a_n\}$

A has no cluster point.

Suppose a_2 is a cluster point.

take $\delta = \min \left\{ \frac{|a_2 - a_1|}{2}, \frac{|a_2 - a_3|}{2} \right\}$, then there is no point of A in $N_\delta(a_2)$.

Finite set has no cluster point.

3. Infinite set $\mathbb{N}$ has no cluster point.

4. $A_n = \left\{ \frac{1}{n} : n \in \mathbb{N} \right\}$ has 0 as a cluster pt.

⑤ $I = [0, 1]$

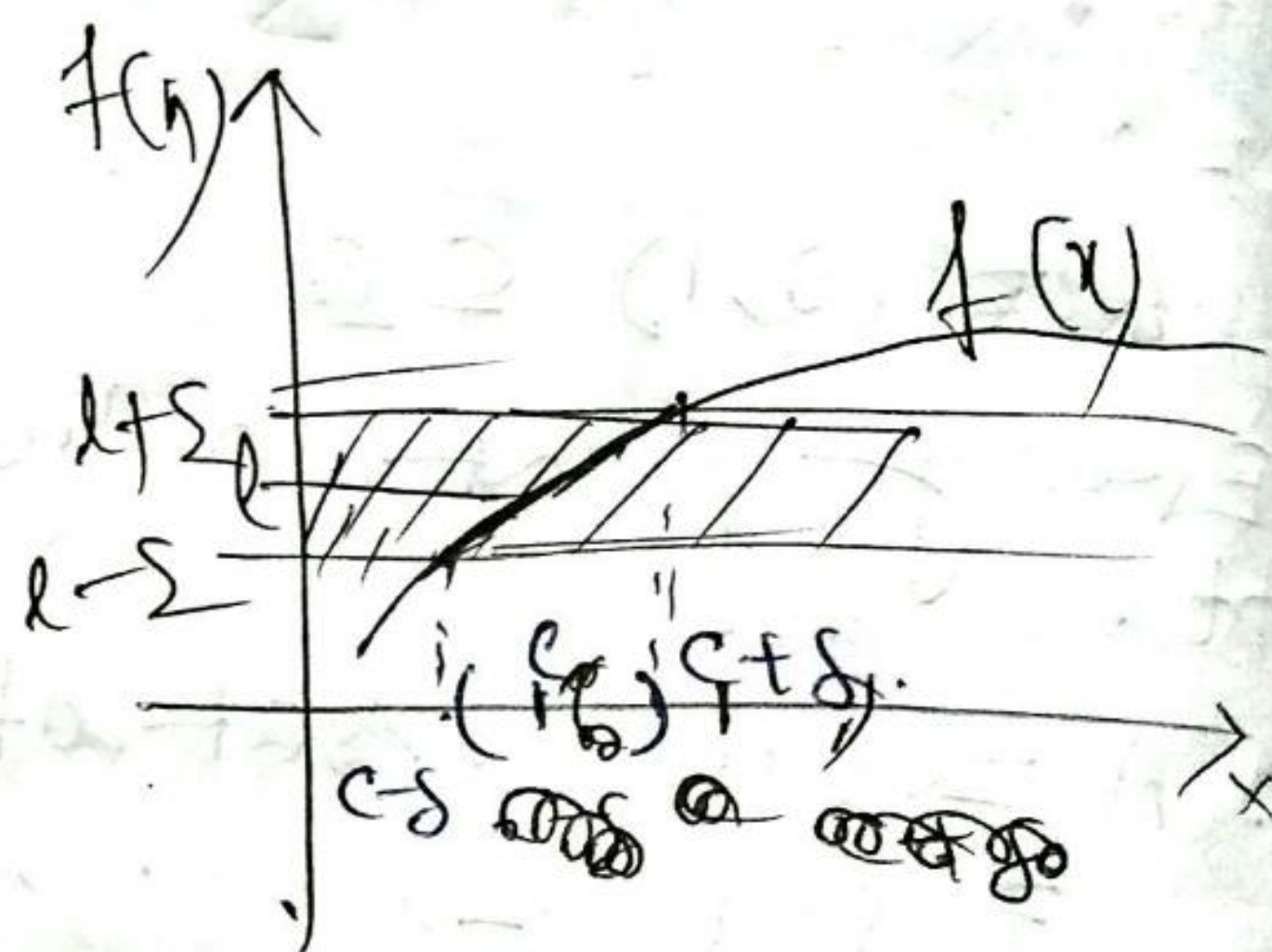
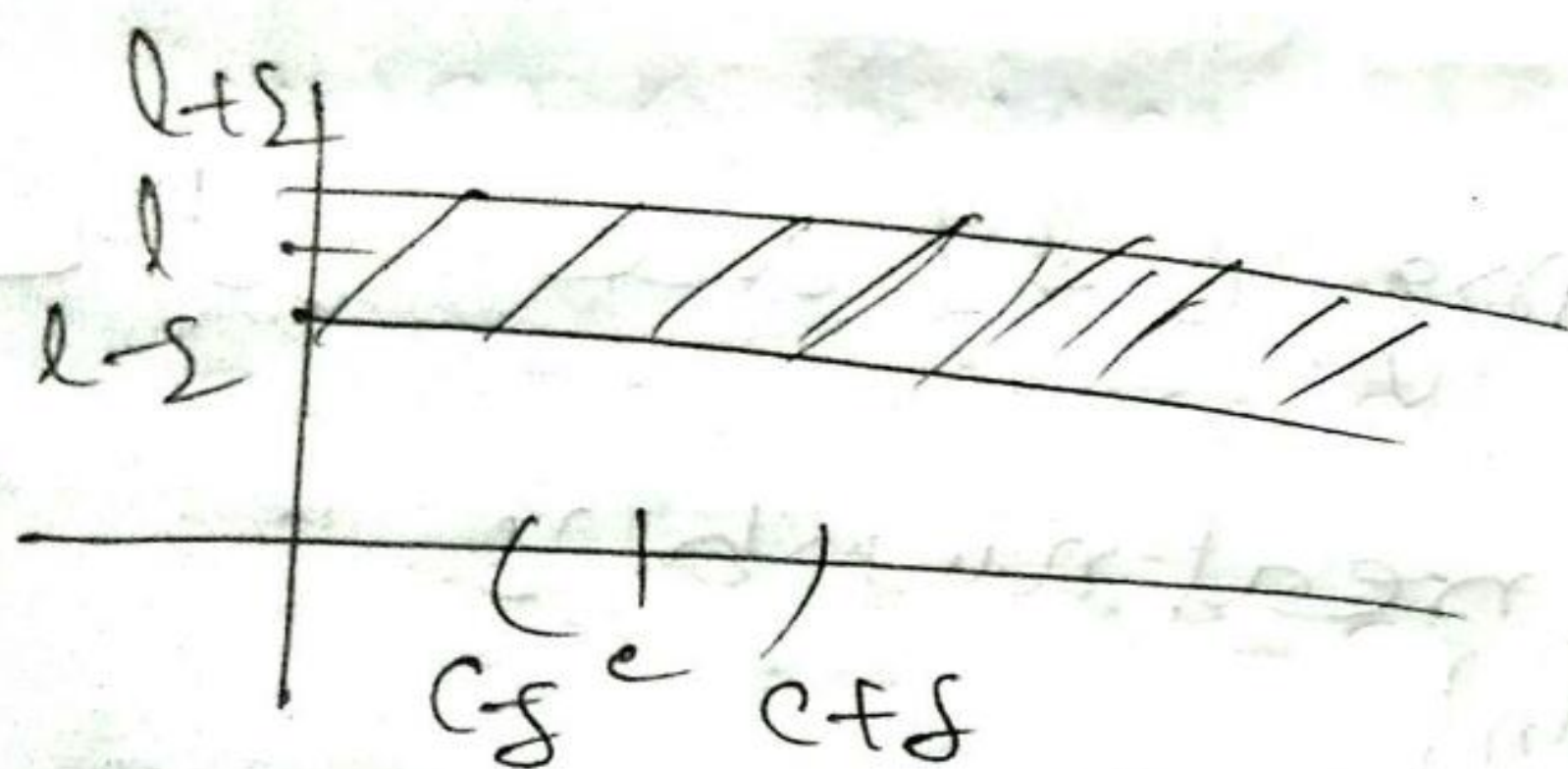
consider $A = I \cap \mathbb{Q}$, $\mathbb{Q}$ is the set of rational numbers.

$A = \text{set of all rational no. lying in } \text{bet}^n [0,1]$

Then every point in I is the cluster pt. of A .

Defⁿ (Limit of function):- Let $A \subseteq \mathbb{R}$, and let c be a cluster point of A .

For a function $f: A \rightarrow \mathbb{R}$, a real number l is said to be a limit of f at c if given any $\epsilon > 0$, there exists a $\delta = \delta(\epsilon) > 0$ such that if $x \in A$ and $0 < |x - c| < \delta$ then $|f(x) - l| < \epsilon$.



It can be easily converted into neighbourhood.

δ -ndd of c , $N_\delta(c) = \{x: |x - c| < \delta\}$

$0 < |x - c| < \delta$ means

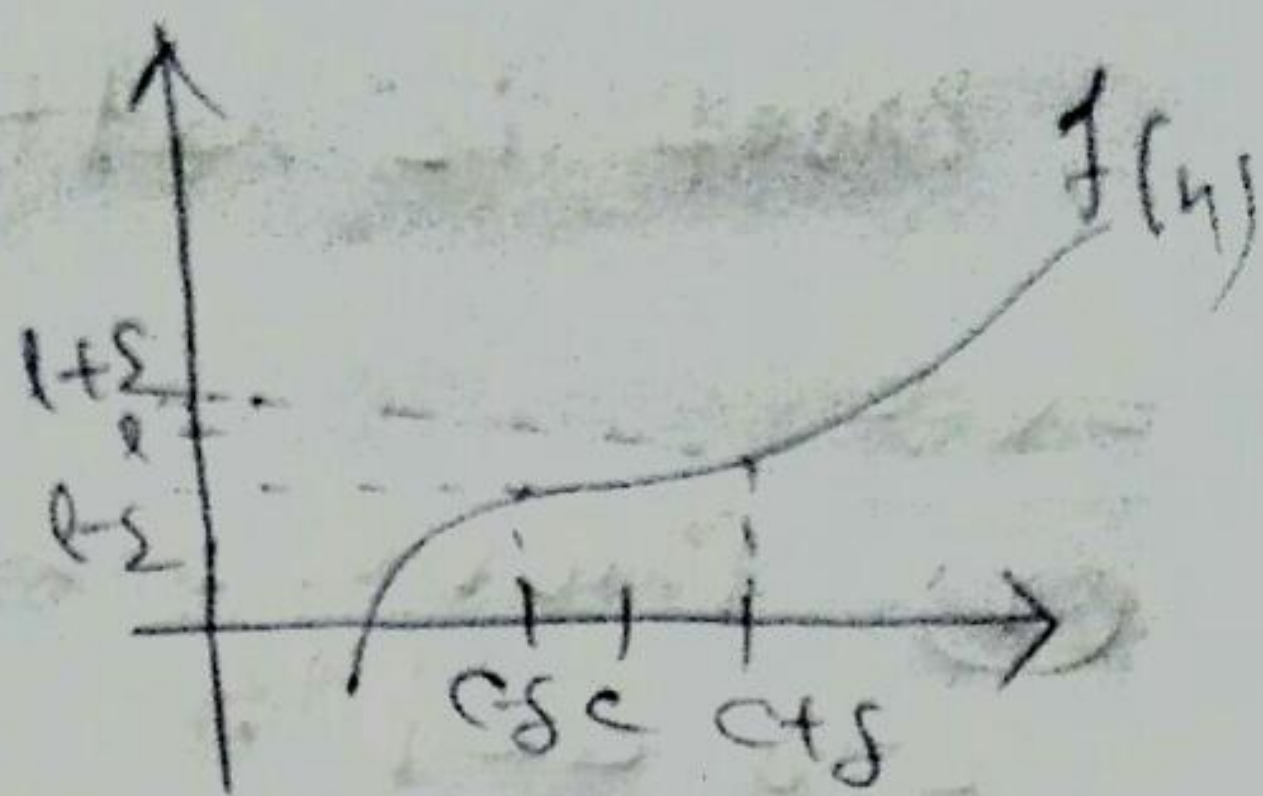
$x \in N_\delta(c)$ but

Similarly, $|f(x) - l| < \epsilon$ means

$x \neq c$.

$f(x) \in N_\epsilon(l)$

If we consider a δ -hbd of l , $N_\delta(l)$, then corresponding to this hbd, we can find a δ -hbd of c , $(c-\delta, c+\delta)$, where c may not be included in it.



Then the limit says that for any arbitrary point x which lies between the interval $(c-\delta, c+\delta)$ and different from c , the corresponding image $f(x)$ will always fall within the interval $(l-\epsilon, l+\epsilon)$.

If it does not happen, then obviously the limit of the fun $f(x)$ when x tends to c will not exist.

It means, there will exist some δ -hbd of l such that whatever the δ -hbd we choose, there will definitely one point whose image will not fall inside $(l-\epsilon, l+\epsilon)$.

Theorem - Let $f: A \rightarrow \mathbb{R}$ and let c be a cluster pt. of A . Then the following statements are equivalent -

(i) $\lim_{x \rightarrow c} f(x) = l$.

(ii) Given any ϵ -hbd of l , i.e., $N_\epsilon(l)$, there exists a δ -hbd $N_\delta(c)$ of c s.t. if $x \neq c$ is any point in $N_\delta(c) \cap A$, then the $f(x) \in N_\epsilon(l)$.

Proof is left for you!

Ex:-

(1) Prove that $\lim_{n \rightarrow \infty} x^2 = c^2$ using ϵ - δ defⁿ.

Here $f(n) = n^2$, $l = c^2$
Given $\epsilon > 0$, we have to identify δ so that the condⁿ given earlier holds.

$$\text{So, } |f(n) - l| = |x^2 - c^2| = |x+c||x-c|$$

Let us take

$$|x-c| < 1$$

$$\text{then } |x| < |c| + 1 \quad \left| \begin{array}{l} 1 > |x-c| \\ \geq |x| - |c| \end{array} \right|$$

$$|x+c| \leq |x| + |c|$$

$$< 2|c| + 1$$

$$\Rightarrow |x| < |c| + 1$$

We are interested to get

$$|f(n) - l| = |x^2 - c^2| = |x+c||x-c| < \epsilon$$

$$|x+c||x-c| < (2|c|+1)|x-c| \text{ when } |x-c| < 1$$

$$\text{Choose } |x-c| < \frac{\epsilon}{2|c|+1}$$

$$\& \delta = \min \left\{ 1, \frac{\epsilon}{2|c|+1} \right\} \quad \text{--- } (**)$$

$$\Rightarrow |x^2 - c^2| < \epsilon, \text{ whenever } 0 < |x-c| < \delta$$

for this particular δ both the results

(*) & (**) will hold together

So, for a given ϵ , we

can identify a

such that

whenever

$$\delta = \min \left\{ 1, \frac{\epsilon}{2|c|+1} \right\}$$

$$0 < |x-c| < \delta, |f(n) - l| < \epsilon$$

$$\Rightarrow \lim_{n \rightarrow \infty} f(n) = l \Rightarrow \lim_{n \rightarrow \infty} n^2 = c^2$$

② To show $\lim_{n \rightarrow 2} \frac{n^3 - 4}{n^2 + 1} = \frac{4}{5}$ using ϵ - δ defn.

Here $f(n) = \frac{n^3 - 4}{n^2 + 1}$, $l = \frac{4}{5}$

Let given $\epsilon > 0$,

$$|f(n) - l| = \left| \frac{n^3 - 4}{n^2 + 1} - \frac{4}{5} \right|$$

$$= \left| \frac{5n^3 - 4n^2 - 24}{5(n^2 + 1)} \right| = \left| \frac{5n^2 + 6n + 12}{5(n^2 + 1)} \right| |n - 2|$$

Let us take

$$|n - 2| < \frac{1}{2} \Rightarrow 1 < n < 3$$

~~first~~ first bound

$$\text{then, } 5n^2 + 6n + 12 < 5 \cdot 9 + 6 \cdot 3 + 12 = 75$$

$$5(n^2 + 1) > 5(1 + 1) = 10$$

$$|f(n) - l| < \frac{75}{10} |n - 2| = \frac{15}{2} |n - 2|$$

$$\text{choose } |n - 2| < \frac{2}{15} \epsilon, \quad \delta = \min \left\{ 1, \frac{2}{15} \epsilon \right\}$$

$$\therefore |f(n) - l| = \left| \frac{n^3 - 4}{n^2 + 1} - \frac{4}{5} \right| < \epsilon \quad \text{provided } |n - 2| < \delta$$

therefore $\lim_{n \rightarrow 2} \frac{n^3 - 4}{n^2 + 1} = \frac{4}{5}$